

Lecture 8.2: Pseudorandom Function Construction

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- Proof: Next Lecture

- Let ℓ and p be prime numbers such that $\ell \mid (p - 1)$

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- Proof: Read on your own

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- Should evaluation of $f_s(x)$ help predict $f_{s'}(x')$?
- Can PRFs be computed by shallow circuits?
[Linial-Mansour-Nisan-94]